

Classify the type of each character

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
<i>S</i>	T	A	T	A	A	T	A	A	T	A	T	A	A	T	A	\$
<i>t</i>	L	S	L	S	S	L	S	S	L	S	L	S	S	L	L	S

S = 1, *L* = 0



②

0 1

0 1 0 1 1 0 1 1

①

1-byte (8 bits) representation

```
void SA_IS(unsigned char *s, int *SA, int n, int K, int cs) {
    int i, j;
    unsigned char *t=(unsigned char *)malloc(n/8+1); // LS-type array in bits
    // Classify the type of each character
    ① tset(n-2, 0); tset(n-1, 1); // the sentinel must be in s1, important!!!
    for(i=n-3; i>=0; i--)
    ② tset(i, (chr(i)<chr(i+1) || (chr(i)==chr(i+1) && tget(i+1)==1))?1:0);
```

$S[i]$ is S-type if $S[i] = S[i+1]$, or $S[i]=S[i+1]$ and $S[i+1, n-1] = S[i+2, n-1]$

```
unsigned char mask[]={0x80, 0x40, 0x20, 0x10, 0x08, 0x04, 0x02, 0x01};
#define tget(i) ( t[(i)/8]&mask[(i)%8] ? 1 : 0 )
#define tset(i, b) t[(i)/8]=(b)?(mask[(i)%8]|t[(i)/8]):((~mask[(i)%8])&t[(i)/8])
#define chr(i) (cs==sizeof(int)?((int*)s)[i]:((unsigned char *)s)[i])
```

In the initial step, *s* is an array of 1-byte characters, but in subsequent steps, *s* can be an array of integers.

Sort all the S-substrings

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
<i>S</i>	T	A	T	A	A	T	A	A	T	A	T	A	A	T	A	\$
<i>t</i>	L	S	L	S	S	L	S	S	L	S	L	S	S	L	L	S
		*		*			*			*		*				*

* = LMS

② Pack LMS from the ends of each bucket.

③ SA:

\$	A									T						
{15}	{-1 -1 -1 -1	{11=09=06=03=01}	{-1 -1 -1 -1 -1 -1}													
@^	^									^						
{15}	{14 -1 -1 -1	{11=09=06=03=01}	{-1 -1 -1 -1 -1 -1}													
^	@	^								^						
{15}	{14 -1 -1 -1	{11=09=06=03=01}	{13 -1 -1 -1 -1 -1}													
^	^		@							^						Some steps are omitted.
{15}	{14 -1 -1 -1	{11=09=06=03=01}	{13 < 10=08=05=02=00}													
^			@													^

①

bkt (size)					
\$	A	C	G	T	
1	9	0	0	6	

②

bkt (end position)					
\$	A	C	G	T	
1	10	10	10	16	

③

bkt (start positions)					
\$	A	C	G	T	
0	1	10	10	10	

This table is used for moving ^ forward.

```

int *bkt = (int *)malloc(sizeof(int)*(K+1)); // bucket array
① getBuckets(s, bkt, n, K, cs, true); // find ends of buckets
   for(i=0; i<n; i++) SA[i]=-1;
② for(i=1; i<n; i++) if(isLMS(i)) SA[--bkt[chr(i)]]+=i;

③ induceSA1(t, SA, s, bkt, n, K, cs, false); // process from the beginning
④ induceSAs(t, SA, s, bkt, n, K, cs, true); // process from the end

```

```

① // find the start or end of each bucket
void getBuckets(unsigned char *s, int *bkt, int n, int K, int cs, bool end) {
    int i, sum=0;
    for(i=0; i<=K; i++) bkt[i]=0; // clear all buckets
    for(i=0; i<n; i++) bkt[chr(i)]++; // compute the size of each bucket
    for(i=0; i<=K; i++) { sum+=bkt[i]; bkt[i]=end ? sum : sum-bkt[i]; }
}

```

end
start

```

② #define isLMS(i) (i>0 && tget(i) && !tget(i-1))

```

i-1	i
L	S
0	1

```

③ void induceSA1(unsigned char *t, int *SA, unsigned char *s, int *bkt,
                int n, int K, int cs, bool end) {
    int i, j;
    getBuckets(s, bkt, n, K, cs, end); // find starts of buckets
    for(i=0; i<n; i++) {
        j=SA[i]-1;
        if(j>=0 && !tget(j)) SA[bkt[chr(j)]++] = j; }
}

```

Sort all the S-substrings

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
<i>S</i>	T	A	T	A	A	T	A	A	T	A	T	A	A	T	A	\$
<i>t</i>	L	S	L	S	S	L	S	S	L	S	L	S	S	L	L	S
		*		*			*			*		*				*

* = LMS

SA: {15} {14 -1 -1 -1 11=09=06=03=01} {13<10=08=05=02=00}

Some steps are omitted.

{15} {14 -1 -1 -1 12<09=07=04=01} {13<10=08=05=02=00}

{15} {14 -1 -1 03<12<09=07=04=01} {13<10=08=05=02=00}

{15} {14 11<06=03<12<09=07=04=01} {13<10=08=05=02=00}

Bkt (ending position)				
\$	A	C	G	T
1	10	10	10	16

This table is used for moving ^ backward.

```

④ void induceSAs(unsigned char *t, int *SA, unsigned char *s, int *bkt,
    int n, int K, int cs, bool end) {
    int i, j;
    getBuckets(s, bkt, n, K, cs, end); // find ends of buckets
    for(i=n-1; i>=0; i--) {
        j=SA[i]-1;
        if(j>=0 && tget(j)) SA[--bkt[chr(j)]] = j;
    }
}
    
```

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
<i>S</i>	T	A	T	A	A	T	A	A	T	A	T	A	A	T	A	\$
<i>t</i>	L	S	L	S	S	L	S	S	L	S	L	S	S	L	L	S
			*	3	*			*	2	*			*	1	*	
			}			}			}			}				
			2			2			3			0				

* = LMS

SA: {15} {14 11<06=03<12<09=07=04=01} {13<10=08=05=02=00}

①

n1=6

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15				
SA:	15	<	11	<	06=	03	<	09=	01		3	2	-1	2	3	1	-1	0	-1	-1

②

③

④

SA: For making the suffix array of 2 1 1 2 0 3 2 2 3 1 0

- ① Pack LMS prefixed into SA according to the temporary ordering, and fill the remaining cells with -1.
- ② Check if neighboring LMS prefixes are equal or not.
- ③ Encode LMS prefixes into non-negative integers according to the ordering.
- ④ Move the encoded non-negative integers to the tail.

```

// compact all the sorted substrings into the first n1 items of SA
// 2*n1 must be not larger than n (proveable)
① int n1=0;
   for(i=0; i<n; i++)
       if(isLMS(SA[i])) SA[n1++]=SA[i];

// find the lexicographic names of all substrings
② for(i=n1; i<n; i++) SA[i]=-1; // init the name array buffer
   int name=0, prev=-1;
   for(i=0; i<n1; i++) {
       int pos=SA[i]; bool diff=false;
       for(int d=0; d<n; d++)
           if(prev==-1 || chr(pos+d)!=chr(prev+d) || tget(pos+d)!=tget(prev+d))
               { diff=true; break; }
           else if(d>0 && (isLMS(pos+d) || isLMS(prev+d))) break;

③   if(diff) { name++; prev=pos; }
       pos=(pos%2==0)?pos/2:(pos-1)/2;
       SA[n1+pos]=name-1;
   }

④ for(i=n-1, j=n-1; i>=n1; i--)
   if(SA[i]>=0) SA[j--]=SA[i];

```

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
<i>S</i>	T	A	T	A	A	T	A	A	T	A	T	A	A	T	A	\$
<i>t</i>	L	S	L	S	S	L	S	S	L	S	L	S	S	L	L	S
		*	3	*			*	2		*	3	*		1		*
		}			}		}			}				}	}	
					2					3				1	0	

* = LMS

SA1																s1						$n-n1=10, n=16, n1=6$											
①																②																	
SA:																→						→											
																5 4 1 2 3 0						3 2 2 3 1 0											
																suffix array of s1						1 3 6 9 11 15											
③																15 11 3 6 9 1						-1 -1 -1 -1 -1 -1 -1 -1 -1 -1											
④																15 11 3 6 9						-1 -1 -1 -1 1						-1 -1 -1 -1 -1 -1 -1 -1 -1 -1					
																15 -1 -1 -1 -1						11 3 6 9 1						-1 -1 -1 -1 -1 -1 -1 -1 -1 -1					

bkt (size)				
\$	A	C	G	T
1	9	0	0	6

bkt (end positions)				
\$	A	C	G	T
1	10	10	10	16

This table is used for moving ^ backward.

- ① Compute the suffix array of **s1**
- ② Decode the nonnegative integers to the original positions of LMS prefixes.
- ③ Translate the suffix array of **s1** into the ordered list of original positions, and fill the remaining cells with -1.
- ④ Move the original positions into proper positions.

```

// stage 2: solve the reduced problem
// recurse if names are not yet unique
int *SA1=SA, *s1=SA+n-n1;
① if(name<n1)
    SA_IS((unsigned char*)s1, SA1, n1, name-1, sizeof(int));
else // generate the suffix array of s1 directly
    for(i=0; i<n1; i++) SA1[s1[i]] = i;

// stage 3: induce the result for the original problem
bkt = (int *)malloc(sizeof(int)*(K+1)); // bucket array

// put all left-most S characters into their buckets
getBuckets(s, bkt, n, K, cs, true); // find ends of buckets
② for(i=1, j=0; i<n; i++)
    if(isLMS(i)) s1[j++]=i; // get p1
③ for(i=0; i<n1; i++) SA1[i]=s1[SA1[i]]; // get index in s
for(i=n1; i<n; i++) SA[i]=-1; // init SA[n1..n-1]
④ for(i=n1-1; i>=0; i--) {
    j=SA[i]; SA[i]=-1;
    SA[--bkt[chr(j)]] = j;
}

```


0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
S T A T A A T A A T A T A A T A \$
t L S L S S L S S L S L S S L L S
 * * * * * * * * = LMS

SA: \$ A T
 { 15 } { -1 -1 -1 -1 < 11 < 03 < 06 < 09 < 01 } { -1 -1 -1 -1 -1 -1 }
 @ ^ ^ ^ ^ ^ ^
 { 15 } { 14 -1 -1 -1 < 11 < 03 < 06 < 09 < 01 } { -1 -1 -1 -1 -1 -1 }
 ^ @ ^ ^ ^ ^ ^
 { 15 } { 14 -1 -1 -1 < 11 < 03 < 06 < 09 < 01 } { 13 -1 -1 -1 -1 -1 }
 ^ ^ @ 途中經過省略 ^ ^
 { 15 } { 14 -1 -1 -1 < 11 < 03 < 06 < 09 < 01 } { 13 < 10 < 02 < 05 < 08 < 00 }
 ^ @ ^ ^

bkt_size				
\$	A	C	G	T
1	9	0	0	6

bkt_start				
\$	A	C	G	T
0	1	10	10	10

This table is used for moving ^ forward.

```
induceSA1(t, SA, s, bkt, n, K, cs, false);
```

- Scan SA from the left by moving “@” on position *i* such that *S*[*i*-1] is L-type, and put position *i* -1 into the open slot labeled with “^” in the bucket of *S*[*i*-1].
- **Proposition:** In each bucket, L-type suffixes come before S-type ones.
- **Proposition:** Suffixes newly added are sorted among each bucket.

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15

S T A T A A T A A T A T A A T A \$

t L S L S S L S S L S L S S L L S

* * * * * * * * = LMS

\$ A T

SA: {15} {14 -1 -1 -1 11<03<06<09<01} {13<10<02<05<08<00}

Some steps are omitted.

{15} {14 -1 -1 -1 12<09<01<04<07} {13<10<02<05<08<00}

{15} {14 -1 -1 06<12<09<01<04<07} {13<10<02<05<08<00}

{15} {14 -1 03<06<12<09<01<04<07} {13<10<02<05<08<00}

{15} {14 11<03<06<12<09<01<04<07} {13<10<02<05<08<00}

bkt_size				
\$	A	C	G	T
1	9	0	0	6

bkt_end				
\$	A	C	G	T
1	10	10	10	16

This table is used for moving ^ backward.

induceSAs(t, SA, s, bkt, n, K, cs, true);

- Scan SA from the right by moving “@” on position i such that $S[i-1]$ is S-type, and put position $i - 1$ into the open slot labeled with “^” in the bucket of $S[i - 1]$.
- **Proposition:** Suffixes newly added are sorted among each bucket.